

A Quantitative Analysis of Algorithms for Energy Efficient Coverage in WSN

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ABSTRACT

In wireless sensor network the Energy Efficient Coverage (EEC) is been a major and important challenge, because sensors work with minimal battery resource in a remote location and it is unfair to change or charge the batteries. An Efficient technique must be implemented as a solution of EEC problem in WSN and it can be achieved by appropriate selection of optimal algorithm. Active research study on EEC problem exposes innovative ideas and solution for coverage issues. This paper discusses different and distinct algorithms have been developed recently for an unstructured WSN to increase network life time and increase coverage among the sensor nodes.

Keywords

Wireless Sensor Network, Energy Efficient Coverage (EEC), jenga inspired optimization algorithm (JOA), TPACO algorithm, ACO, ACB-SA algorithm, optimal lifetime enhancement

1. INTRODUCTION

Wireless sensor network often encompasses with sensors, processors, transceiver and battery. In WSN, sensor collects data about the environment (surroundings) or from other sensor nodes and processes data finally, it will be communicated with other sensor nodes and base station. Due to the growth in advanced embedded systems and MEMS concept sensors performs better and high performance sensors are available at cheaper cost with low power batteries. Usually WSN implementation held at the environment changes and to be monitored.

Coverage is the fundamental problem in any wireless network especially WSNs. In WSN the coverage is not only based on sensor nodes, it depends on several parameters such as sensor speed, memory involved, algorithms and scheduling methods followed etc., In general WSNs work with low power, short life span and tiny sensors. These sensors are more often deployed in remote or inaccessible locations. Hence, it is difficult to conserve network energy to prolong lifetime of WSNs. Network performance definitely affects the coverage and considered measure of Quality of Service (QoS) in WSN. Some scheduling and planning algorithms reduces energy consumption to increase life time. Activity scheduling problem deals with different coverage issues. It can be classified as target coverage, area coverage, barrier coverage and patrol coverage. Three centralized and one distributed energy efficient coverage algorithms discussed in [15]. For full area problems BEFAC algorithm was introduced and it uses minimum no. of sensors [5] and in [8] multi objective genetic algorithm was proposed for full coverage issues. To improve network life in WSN, scheduling algorithms and methods were followed based on ant colony system [7].

Integrated coverage analysis coverage configuration protocol and relationship between coverage and connectivity [5] [6]. Reliable and energy efficient coverage for WSN approach ensures the successful coverage of targets [14]. ACO based energy saving routing reformulates energy consumption in WSN. To maximize energy levels by putting more number of sensors in sleep CDS based coverage protocols discussed in [16].

The main objective of this paper is to explore different types of algorithms proposed by researchers recently for optimal energy efficient coverage issue in WSN. Another goal is to discourse the algorithm to be used in real time implementation of WSN for the better battery life as well as network lifetime. Researchers focus on coverage issue in WSN to enhance lifetime and reduce processing time to improve quality of service between sensor nodes. In next section energy efficient coverage problem is defined and the assumptions made for ACO, TPACO, ACB-SA and Jenga inspired optimal algorithm are discussed.

2. ENERGY EFFICIENT COVERAGE

Given a monitored region A, a set S of sensors ($S = \{s_1, s_2, s_3, \dots, s_{N_s}\}$) a set P of POI $P = \{p_1, p_2, p_3, \dots, p_{N_i}\}$ and a cost of C_i of the sensor i ($i = 1, 2, 3, \dots, N_s$) find a set C of subsets $C(t_s)$ ($t_s = 1, 2, 3, \dots, T_s$) of sensors that cover all of the POI in region A with a minimum total cost for a time slot until the WSN fails to cover any of the POIs in a region A. The T_s is the lifetime of WSN [3] [1] & [4].

In a real time deployed sensors detect events (physical parameters) like temperature, humidity that occur at a point of interest by measuring the received energy. The received signal strength is greatly reduced exponentially when the distance between POI and sensor increases. The formulation model was introduced in [9]

$$\lambda_i(j) = \begin{cases} 0, & \text{if } d_{ij} > r_u \\ e^{-a(d_{ij}-r_s)^m}, & \text{if } r_s < d_{ij} \leq r_u \\ 1, & \text{if } d_{ij} \leq r_s \end{cases} \quad (1)$$

λ_i indicates the probability of detection of events at POI j by sensor i, and d_{ij} is the Euclidean distance between sensor i and POI j. In equation (1) m and a are variable (delay factors). From the above equation the following points could be observed.

1. If the distance of POI j is in the range r_s , event at POI j are definitely detected.
2. Whenever the distance is greater than r_s , the probability of detection decreases exponentially.
3. If the distance is greater than r_u and the received signal strength is smaller than that of noise produces the probability zero.

Some assumptions made for the solution of EEC problem for the following algorithms based on the results of [6] [9].

- The communication range r_c of the device at least twice its sensing range.
- The positions of the sensors and POI are well defined.
- The sensors do not consume energy when it is in inactive mode. Finally, cost of an active sensor i is predicted, $C_i = k_{ER,i}$ (2)
 E_R is the residual energy, k is the constant based on sensor characteristics.

The probability of detection $\lambda(j)$ is larger than ϵ and it is defined by the user.

$$\lambda(j) = 1 - \prod_{s_i \in S_j} (1 - \lambda_i(j)) \geq \epsilon \quad (3)$$

Where, S_j is a set of sensors that covers the POI j $S_i \subset S$. With reference to [9], the EEC problem can be expressed as,

$$\text{Minimize, } \sum_{i=1}^{N_s} C_i x_i \quad (4)$$

Subject to

$$\forall p_j \in P, \sum_{i=1}^{N_s} \alpha_{ij} x_i \geq \delta, x_i \in \{0,1\} \quad (5)$$

δ is a variable and it is the smallest number of sensors. α_{ij} is a function of $\lambda_i(j)$ and $\alpha_{ij} = -\ln(1 - \lambda_i(j))$.

2.1 Conventional ACO Algorithm

Ant colony optimization algorithm is based on the behavioral approach of ants. Ant search for food and it will move in and around while moving, it deposits pheromones and these pheromones helps and lead remaining ants for path identification. The pheromones evaporate with time. Then ants establishes new path and will lay pheromones. This method is suitable to find shortest path in WSN to transfer content of information in between sensor node to other nodes or base station. This is energy efficient and it was applied first in TSP [10]. ACO algorithm for the TSP, concerns N cities and M ants and the mathematical expression as follows for the transition probability from city i to city j for the k^{th} ant.

$$P_{ij}^k(t) = \begin{cases} \frac{[\tau_{ij}(t)]^\alpha [\eta_{ij}]^\beta}{\sum_{k \in allowed_k} [\tau_{ik}(t)]^\alpha [\eta_{ik}]^\beta}, & \text{if } j \in allowed_k \\ 0, & \text{otherwise} \end{cases} \quad (6)$$

Where, $allowed_k$ is the remainder cities to be covered. α and β are the constants and it is relative influence factor of pheromone. For every tour or travel of ant the pheromones amount will be updated based on local pheromone decay parameter ρ and π (0, 1) and the formula can be written,

$$\tau_{ij}(t+N) = (1-\rho)\tau_{ij}(t) + \Delta\tau_{ij} \quad (7)$$

If τ is the initial update and added pheromone amount calculated for $t+N$ point.

$$\Delta\tau_{ij} = \sum_{k=1}^M \Delta\tau_{ij}^k \quad (8)$$

Where $\Delta\tau_{ij}$, is the amount per unit length of pheromone trail on edge (i,j) by the k^{th} ant between t and time $t+N$ and is given as follows:

$$\Delta\tau_{ij}^k = \begin{cases} \frac{Q}{L_k}, & \text{if } k^{\text{th}} \text{ ant uses edge}(i,j) \text{ in its tour} \\ 0, & \text{otherwise} \end{cases} \quad (9)$$

Where, Q is the constant, L_k is length of the path established by k^{th} ant (sensor). This process continues up to find optimal path for EEC.

2.2 Ant Colony Optimization with Three Pheromones

This algorithm refers conventional ant colony system algorithm. But, it improves lifetime of network. TPACO algorithm considers three pheromones unlike ACO uses one pheromone [4]. In three pheromones, one is local pheromone and used to organize coverage with sensors. Remaining two pheromones are called global pheromones. One of the global pheromone customizes the number of needed active sensors for each POI and another is used to form a sensor set that is based on former pheromone.

The TPACO algorithm did not consider the parameters of α and β which are used as the function in conventional ACO. As the first step of TPACO algorithm position information of sensors and POIs are collected and stored as matrix. The matrix helps to initialize global pheromone field T_{AS} per POI and for a time slot ($t_s = 1, 2, 3.. T_s$). Whenever the fewer number of active sensors used it definitely enhances the life time of WSN because most of the deployed sensors will be in sleep mode. T_{NoAS} is initialized by Gaussian function.

$$T_{NoAS, jm}(0) = \begin{cases} \frac{1}{\sigma\sqrt{2\pi}} e^{-(m-\mu_j)/2\sigma^2}, & \text{if } m \leq n_j \\ 0, & \text{otherwise} \end{cases} \quad (10)$$

Where, n_j is the number of sensors covered at POI j . $m = (1, 2, n_j)$ σ is constant μ_j is mean. Initially mean is zero and when the ant fails to organize the sensor set mean increases. Repeated failure causes unpredictability of number of sensor at POI j .

As the next stage, ant (sensor) k plays roulette wheel selection probability of sensor $S_{i,j}^k$ for ant k .

$$P_{AS}(S_{i,j}^k) = \begin{cases} \frac{\tau_{AS,ji}(t) + \tau_{SS,i}^k(t)}{\sum_{m \in allowed_m} (\tau_{AS,jm}(t) + \tau_{SS,m}^k(t))}, & \text{if } i \in allowed_m \\ 0, & \text{otherwise} \end{cases} \quad (11)$$

Where, $T_{AS, ji}(t)$ is the global variable. $\tau_{SS,i}^k(t)$ is the local variable. $allowed_m$ is the set of remaining sensors. The local pheromone is updated each and every time ant k travels for POI j . For every time slot t_s , the updated information covers the POIs. This process continues up to the failure of network to cover any POI as per basic equation (5).

2.3 V. ACB – SA Algorithm

To improve the lifetime of network ACB SA was introduced in [1]. It is the modified algorithm of conventional ACO. This algorithm solves the problem of imaginary ant behavioral characteristics based on the previous construction graph. In conventional ACO if POI is not covered by sensor (ant k) it

will refer next sensor (ant k+1). In ACB SA parameters α and β are not considered. The algorithm begins with pheromone field τ_s .

$$Ts.i(t_s) = E_{r,i}(t_s) \cdot I_{cover\ x,i}(t_s) \quad (12)$$

Where, $E_{r,i}(t_s)$ is the residual energy of sensor i at time slot t_s . I_{cover} is the collection of information about covered POI.

$$p_s(i) = \frac{\tau_{s,i}(n_C)}{\sum_{i=1}^{N_s} \tau_s(n_C)} \quad (13)$$

$P_s(i)$ is the sensor selection probability. After completion of covering all the POIs, next sensor (ant k+1) will be selected.

2.4 Jenga Optimal Algorithm

Jenga Inspired Optimal Algorithm (JOA) was proposed [13] based on Jenga board game. In this algorithm, probabilistic sensor detection model was introduced which solves the EEC problem. The positions information are collected from sensors and POIs and stored as a matrix. This matrix includes residual energy of sensors. As a next stage set C is initialized to store subsets of sensors in each time slot and another variable is introduced to store the life time of WSN for EEC. The user parameters of jenga algorithm i.e. n_p (number of players) and n_T (turns) and E for the smallest probability of detection of POI are defined. Once algorithm begins to function, for every timeslot t_s the scoreboard Ω_s is established. The value of element is calculated by

$$\Omega_{s,i}(t_s) = E_{R,i}(t_s) I_{cover,i}(t_s) \quad (14)$$

Where, $I_{cover,i}(t_s)$ is the no. of POIS covered by sensor i .

In Jenga algorithm nodes take turns deactivating a sensor from a network of active network sensors. It is considered by probability selection model and is given by,

$$p_s(i) = \frac{\Omega_{s,i}(n_T)}{\sum \Omega_s(n_T)} \quad (15)$$

Player k try to remove sensor within turn n_T which is based on equation (5). If k fails next sensor $K+1$ will take turn and it continues/ this process will be updated with set C and the optimal cover set is considered. If WSN is unable to cover POI, another node takes a turn. This process continues to satisfy POIs. The JOA was implemented in matlab for different numbers of sensors and POIs to calculate life time and computation time

3. RESULTS & INTERPRETATIONS

This section is fully dedicated for analyze the simulation results of algorithms studied earlier in this paper. The main objective of this paper is to discuss the network life time, complexity and the computational time of algorithms used to solve EEC problem in WSN. From the simulation results of [1], [4] & [13] table 1, table 2 and table 3 are made. Table 1 and fig.1 describes the average network lifetime of different algorithms suited in WSN. This average is calculated after performing simulation about 30 times and concerns different scenarios. In scenario 1 considers 100 sensors and 10 POI, scenario 2 applies 150 sensors and 10 POI and in scenario 3 200 sensors taken for 10 POIs.

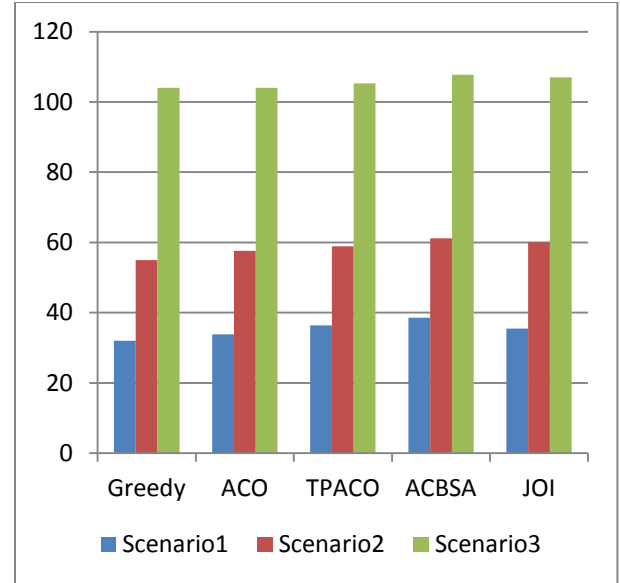


Fig 1: Average Network Lifetime for Algorithms in Seconds

Table 1. Average Network Lifetime for Algorithms in Seconds

	Greedy	ACO	TPACO	ACBSA	JOA
Scenario1	32	33.8	36.4	38.6	35.5
Scenario2	55	57.6	58.9	61.2	60.1
Scenario3	104	104	105.3	107.7	107

The reason hidden in to analyze complexity and computational time in WSN is, WSN is the combination of sensors, transceivers, processors; it is the co design issue, whether to choose long battery life or high processing speed. Actually the complexity of the network system algorithm is calculated based on the number of sensors deployed, no. of agents, no. of active sensors and no. of iteration process using Big O notation. Complexity of an algorithm refers how much repetitive process going on to find result (optimal coverage) for one time slot. The need to analyze complexity is it affects the performance of the system. Table 2 and fig.2 explores maximum complexity of algorithm for one time slot.

Table 2. Complexity of Algorithms

Algorithms	Big - O notations	Complexity (maximum no of iterations for one timeslot)
Greedy	N^2	200X30
ACO	N^3	100X500X30
TPACO	N^4	30X300X30X n_j^k
ACBSA	N^3	50X200X30
JOA	N^3	100X100 X 30

n_j^k refers maximum value of 200

Figure 2 shows the average computational time for different algorithms discussed and it uses only 10 POIs and 100, 150 and 200 sensors respectively for scenarios.

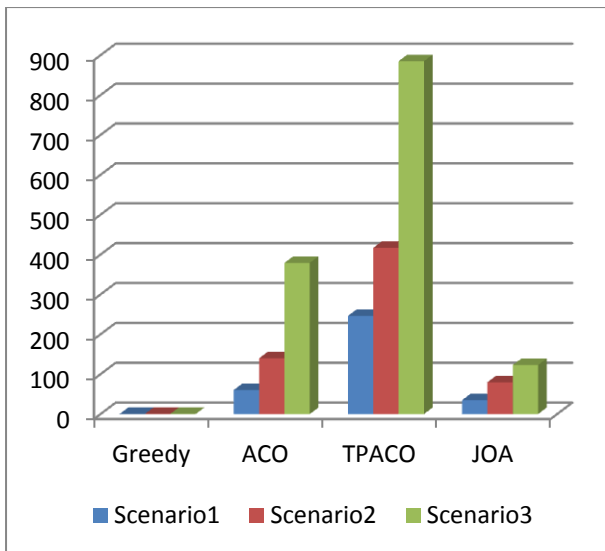


Fig 2: Computational Time of Algorithms in Sec

Figure 3 shows the average computational time for various algorithms and it takes 20 POIs and 100, 150 and 200 sensors for scenarios 4, 5 and 6 respectively.

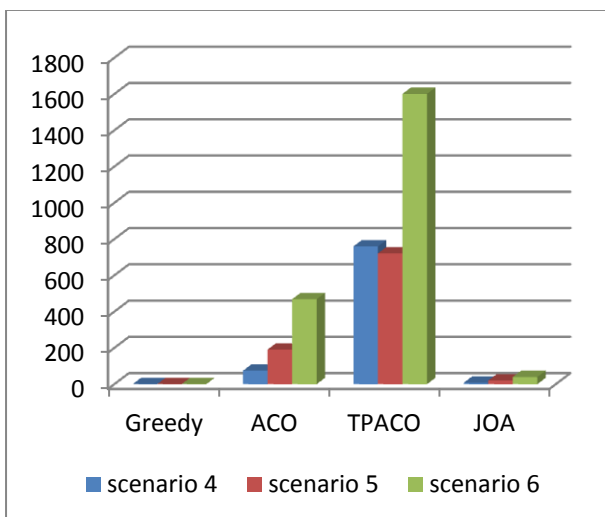


Fig 3: Computational Time of Algorithms in Sec

Figure 4 shows the average computational time for different algorithms and it has 30 POIs and 100, 150 and 200 sensors for scenarios 7,8 and 9 respectively.

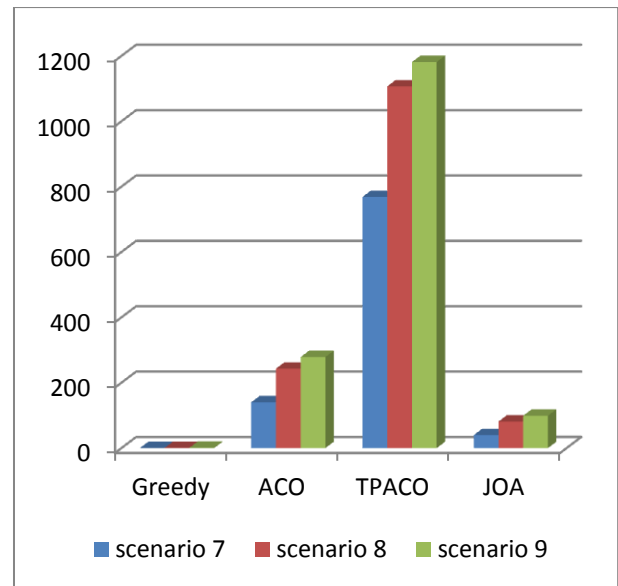


Fig 4: Computational Time of Algorithms in Sec

4. CONCLUSIONS

In this paper, different optimal algorithms used recently for EEC problem in WSN such as, ACO, TPACO, ACBSA and JOA are discussed. As an outcome, each algorithm has its unique characteristics. In greedy method, the computational time and complexity are low but network lifetime is poor. Greedy algorithm uses mathematical calculation only. Ant colony optimization algorithm is a good solution for finding shortest path between the sensor nodes but with respect to alpha and beta the entire performance of the network system will vary. Three pheromones ACO is a new approach and it removes alpha and Beta instead local & global pheromones are initialized to update activity of sensor to the base station. Probably, for small level systems in WSN having less than 200 sensors may apply TPACO algorithm for EEC problem. The ACBSA is more realistic approach, which provide longer lifetime and parameters selection was made by probabilistic sensor detection model. Result found that ACBSA is better algorithm than TPACO and ACO in network lifetime perspective. Contrary, to greedy algorithm the complexity is very high. Finally, JOI algorithm widened the view of optimal algorithm development. JOI algorithm performs based on the Jenga board game. In that, computational time is nominal and the network lifetime also better than TPACO, ACO and Greedy algorithms. As a conclusion, due to active research in WSN especially for EEC problem in future, optimal algorithms and solutions are expected and it is inevitable.

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