A High Spectral Efficient Non-Binary TTCM-Assisted G2 STBC-OFDM for 4G Systems

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ABSTRACT

This paper deals with the MIMO-OFDM technique that candidate with the fourth generation (4G) of the wireless communication systems, this technique can provide high data rate transmission without increasing transmit power and expanding bandwidth, also it can efficiently use space resources and has a bright future. This paper presents the channel coding assisted STBC-OFDM systems, and employs the Coded Modulation techniques (CM), since the signal bandwidth available for wireless communications is limited. Furthermore, a non-binary Turbo Trellis Coded Modulation (TTCM) decoder-based multidimensional MAP algorithmassisted G2 STBC-OFDM is design and implemented. The idea of Non-binary codes has been extended for symbols defined over rings of integers, which outperform binary codes with only a small increase in decoding complexity. The simulation results show that the performance of the TTCM decoding algorithm outperforms the binary decoding methods and suitable to deal with error control coding of the STBC-OFDM schemes.

Keywords

Turbo codes, TTCM, non-binary error correcting codes, Rings of integers, MIMO, OFDM MIMO-OFDM, STBC, 4G.

1. INTRODUCTION

Increasing market expectations for 3G (third generation) systems show a great demand for a wider range of services spanning from voice to high-rate data services required for supporting mobile multimedia communications. On the other hand, the available radio spectrum is limited and the communication capacity needs cannot be met without a significant increase in spectral efficiency.

4G technology is the latest technology after the 3G systems. These 4G Communication systems still solve the remaining problems of 3G systems. 4G is characterized by three words: ubiquitous, mobile, and broadband [1]. The main key features of 4G are receiving a large volume of information, data, pictures, video and so on. A MIMO system is capable of exploiting transmitter and receiver diversity, hence maintaining reliable communications.

As a key building block of next-generation wireless communication systems, MIMOs are capable of supporting significantly higher data rates than the Universal Mobile Telecommunications System (UMTS) and the High-Speed Downlink Packet Access (HSDPA) based 3G networks [2].

To overcome a multipath-fading environment with low complexity and to achieve wireless broadband multimedia communication systems (WBMCS), the orthogonal frequency-division multiplexing (OFDM) transmission scheme is employed [3]. OFDM is one of the applications of a parallel-data-transmission scheme, which reduces the influence of multipath fading and makes complex equalizers unnecessary.

Advances in coding, such as turbo [4] and low density parity check codes [5], made it feasible to approach the Shannon capacity limit [6] in systems with a single antenna link. Significant further advances in spectral efficiency are available through increasing the number of antennas at both the transmitter and the receiver [7, 8, 9].

Further performance gains can be achieved by using nonbinary codes in the coded modulation scheme, but with an increase in the decoding complexity [10].

J. Kumari J, 2011 [11] analyzed the performance of MIMO-OFDM system employing Quadrature Amplitude Modulation (QAM) in Rayleigh fading channel, since the use of multiple antennas at both ends of a wireless link (multiple-input multiple output (MIMO) technology) has recently been demonstrated to have the potential of achieving extraordinary data rates. Orthogonal frequency division multiplexing (OFDM) significantly reduces receiver complexity in wireless broadband systems. The use of MIMO technology in combination with OFDM, i.e., MIMO-OFDM therefore seems to be an attractive solution for future broadband wireless systems.

L. Hanzo et. al., 2011 [12] presented that the CM-assisted STBC schemes were found to improve significantly the system's achievable performance. Furthermore, the TTCM-STBC concatenated scheme was observed to give the best performance among all the CM-STBC concatenated schemes. Also, the TTCM-assisted G2 coded OFDM scheme gives a better performance than the LDPC-assisted G2 coded OFDM scheme. Furthermore, in the context of the achievable coding gain versus complexity performance, it was found that the TTCM-assisted schemes are capable of achieving higher coding gains in the relatively high-complexity range, than the LDPC-assisted candidate schemes.

2. OFDM- MIMO, and MIMO-OFDM

The principle of multi-carrier transmission is to convert a serial high rate data stream on to multiple parallel low rate sub-streams. Each sub-stream is modulated on another sub-carrier. Since the symbol rate on each sub-carrier is much less than the initial serial data symbol rate, the effects of delay spread, i.e. ISI, significantly decrease, reducing the complexity of the equalizer. OFDM is a low complex technique used to modulate multiple sub-carriers efficiently by using digital signal processing [13, 14, 15, 16, 17]. An important design goal for a multi-carrier transmission scheme based on OFDM in a mobile radio channel is that the channel

and that fading per sub-channel can be considered as flat. As conventional methods like using more bandwidth or higher order modulation types are limited, new methods of using the transmission channel have to be used. Multiple antenna systems (Multiple Input, Multiple Output– MIMO) give a significant enhancement to data rate and channel capacity. A MIMO system employs multiple transmitter and receiver antennas for delivering parallel data streams, as illustrated in Figure (1). Since the information is transmitted through different paths, a MIMO system is capable of exploiting transmitter and receiver diversity, hence maintaining reliable communications.



Fig 1: A MIMO System.

As compared with Single-Input, Single-Output (SISO) systems, the most significant advantages of MIMO systems are; a significant increase of both the system's capacity and spectral efficiency, and dramatic reduction of the effects of fading due to the increased diversity.

The quality of a wireless link can be described by three basic parameters, namely the transmission rate, the transmission range and the transmission reliability. Conventionally, the transmission rate may be increased by reducing the transmission range and reliability. By contrast, the transmission range may be extended at the cost of a lower transmission rate and reliability, while the transmission reliability may be improved by reducing the transmission rate and range [18].

However, with the advent of MIMO-assisted OFDM systems, the above-mentioned three parameters may be simultaneously improved. Initial field tests of broadband wireless MIMO-OFDM communication systems have shown that an increased capacity, coverage and reliability are achievable with the aid of MIMO techniques [19].

A generic MIMO-OFDM system employing K orthogonal frequency-domain subcarriers and having mt and nr transmit and receive antennas, respectively.

Firstly, the OFDM modulation technique is capable of coping with the highly frequency-selective, time-variant channel characteristics associated with mobile wireless communication channels, while possessing a high grade of structural flexibility for exploiting the beneficial properties of MIMO architectures.

The family of space-time signal processing methods, which allow for the efficient implementation of communication systems employing MIMO architectures, is commonly referred to as smart antennas.

The schematic of a typical MIMO-OFDM system's physical layer is depicted in Figure (2). The transmitter of the MIMO-OFDM system considered is typically constituted by the encoder and modulator, generating a set of mt complexvalued baseband time-domain signals. The modulated baseband signals are then processed in parallel.

Specifically, they are over sampled and shaped using a Nyquist filter. The resultant over sampled signals are then converted into an analogue pass band signal using a bank of D/A converters and up converted to the RF band.

At the receiver side of the MIMO-OFDM transceiver, the inverse process takes place, where the set of received RF signals associated with the nr receive antenna elements is amplified by the RF amplifier and down converted to an intermediate-frequency passband.

The resultant passband signals are then sampled by a bank of A/D converters, down converted to the baseband, filtered by a matched Nyquist filter and finally decimated, in order to produce a set of discrete complex-valued baseband signals. The resultant set of discrete signals is processed by the corresponding demodulator and decoder module seen in Figure (2) where the transmitted information-carrying symbols are detected.



Fig 2: Schematic of a typical MIMO-OFDM system's physical layer

The discrete frequency-domain model of the MIMO-OFDM system is illustrated in Figure (2) may be characterized as a generalization of the Single-Input Single-Output (SISO) case:

$$y_{i}[n,k] = \sum_{i=1}^{m_{t}} H_{ij}[n,k] x_{j}[n,k] + w_{i}[n,k] \quad (1)$$

where n =0, 1,... and k =0,...,K – 1 are the OFDM symbol and subcarrier indices, respectively, while yi[n, k], xj [n, k] and wi[n, k] denote the symbol received at the ith receive antenna, the symbol transmitted from the jth transmit antenna and the Gaussian noise sample encountered at the ith receive antenna, respectively. Furthermore, Hij [n, k] represents the complexvalued CTF coefficient associated with the propagation link connecting the jth transmit and ith receive antennas at the kth OFDM subcarrier and time instance n. The MIMO-OFDM system model described by (2) can be interpreted as the per-OFDM- subcarrier vector expression of

 $\mathbf{y}[\mathbf{n}, \mathbf{k}] = \mathbf{H}[\mathbf{n}, \mathbf{k}]\mathbf{X}[\mathbf{n}, \mathbf{k}] + \mathbf{W}[\mathbf{n}, \mathbf{k}] \quad (2)$

3. Alamouti's G2 STBC

The system contains two transmitter antennas and one receiver antenna, a generic STBC is defined by an $(n \times p)$ -dimensional transmission matrix G, where the entries of the matrix G are linear combinations of the k input symbols x1, x2, ..., xk and their conjugates. Each symbol xi (i =1,..., k) conveys b original information bits according to the relevant signal constellation that has M =2b constellation points, and

hence can be regarded as information symbols. Thus, $(k \times b)$ input bits are conveyed by each $(n \times p)$ block.

The G2 transmission matrix can be derived in the form of [20].

$$\begin{bmatrix} x_1 & x_2 \\ -\overline{x}_2 & \overline{x}_1 \end{bmatrix}$$
(3)

Since there are k = 2 input symbols, namely x1 and x2, the code rate of G2 is R = k/n = 1.

Two algorithms are widely used for decoding STBCs. The maximum likelihood (ML) decoding algorithm generates hard-decision outputs, while the Maximum-A-Posteriori (MAP) decoding algorithm is capable of providing soft outputs, which readily lend themselves to channel coding for achieving further performance improvements.

3.1 Maximum-A-Posteriori Decoding

Bauch [21] presented a simple symbol-by-symbol MAP algorithm for decoding STBCs. According to [21], the a posteriori probability of each information symbol xi (i =1,...,k) is given by

$$\ln P(x_i|y_1, y_2, ..., y_q) = const + \ln P(y_1, y_2, ..., y_q|x_i) + \ln P(x_i)$$
(4)

Where $[y_1^j, y_2^j, \dots, y_n^j]$ (j =1,...,q) represents the received signal vector at receiver j during the period spanning from time slot 1 to time slot n, for the STBC G2 (k = 2, n = 2) the a posteriori probabilities are given by

$$\ln P(x_1 | y_1, y_2, ..., y_q) = const - \frac{1}{2\sigma^2} \{ \left| \left[\sum_{j=1}^q (y_1^j \, \overline{h}_{1,j} + y_2 / h_2, j) - x_1 2 + j = 1 q t = 12 h t_j / 2 - 1 x_1 2 + \ln P(x_1) \right]$$

$$\ln P(x_2 | y_1, y_2, ..., y_q) = const - \frac{1}{2\sigma^2} \{ \left| \left[\sum_{j=1}^q (y_1^j \, \overline{h}_{2,j} + y_1^j) - y_1 \right] \right| = const - \frac{1}{2\sigma^2} \{ \left| \sum_{j=1}^q (y_1^j \, \overline{h}_{2,j} + y_1^j) - y_1 \right| = const - \frac{1}{2\sigma^2} \{ \left| \sum_{j=1}^q (y_1^j \, \overline{h}_{2,j} + y_1^j) - y_1 \right| = const - \frac{1}{2\sigma^2} \{ \left| \sum_{j=1}^q (y_1^j \, \overline{h}_{2,j} + y_1^j) - y_1 \right| = const - \frac{1}{2\sigma^2} \{ \left| \sum_{j=1}^q (y_1^j \, \overline{h}_{2,j} + y_1^j) - y_1 \right| = const - \frac{1}{2\sigma^2} \{ \left| \sum_{j=1}^q (y_1^j \, \overline{h}_{2,j} + y_1^j) - y_1 \right| = const - \frac{1}{2\sigma^2} \{ \left| \sum_{j=1}^q (y_1^j \, \overline{h}_{2,j} + y_1^j) - y_1 \right| = const - \frac{1}{2\sigma^2} \{ \left| \sum_{j=1}^q (y_1^j \, \overline{h}_{2,j} + y_1^j) - y_1 \right| = const - \frac{1}{2\sigma^2} \{ \left| \sum_{j=1}^q (y_1^j \, \overline{h}_{2,j} + y_1^j) - y_1 \right| = const - \frac{1}{2\sigma^2} \{ \left| \sum_{j=1}^q (y_1^j \, \overline{h}_{2,j} + y_1^j) - y_1 \right| = const - \frac{1}{2\sigma^2} \{ \left| \sum_{j=1}^q (y_1^j \, \overline{h}_{2,j} + y_1^j) - y_1 \right| = const - \frac{1}{2\sigma^2} \{ \left| \sum_{j=1}^q (y_1^j \, \overline{h}_{2,j} + y_1^j) - y_1 \right| = const - \frac{1}{2\sigma^2} \{ \left| \sum_{j=1}^q (y_1^j \, \overline{h}_{2,j} + y_1^j) - y_1 \right| = const - \frac{1}{2\sigma^2} \{ \left| \sum_{j=1}^q (y_1^j \, \overline{h}_{2,j} + y_1^j) - y_1 \right| = const - \frac{1}{2\sigma^2} \{ \left| \sum_{j=1}^q (y_1^j \, \overline{h}_{2,j} + y_1^j) - y_1 \right| = const - \frac{1}{2\sigma^2} \{ \left| \sum_{j=1}^q (y_1^j \, \overline{h}_{2,j} + y_1^j) - y_1^j \right| = const - \frac{1}{2\sigma^2} \{ \left| \sum_{j=1}^q (y_1^j \, \overline{h}_{2,j} + y_1^j) - y_1^j \right| = const - \frac{1}{2\sigma^2} \{ \left| \sum_{j=1}^q (y_1^j \, \overline{h}_{2,j} + y_1^j) - y_1^j \right| = const - \frac{1}{2\sigma^2} \{ \left| \sum_{j=1}^q (y_1^j \, \overline{h}_{2,j} + y_1^j) - y_1^j \right| = const - \frac{1}{2\sigma^2} \{ \left| \sum_{j=1}^q (y_1^j \, \overline{h}_{2,j} + y_1^j) - y_1^j \right| = const - \frac{1}{2\sigma^2} \{ \left| \sum_{j=1}^q (y_1^j \, \overline{h}_{2,j} + y_1^j) - y_1^j \right| = const - \frac{1}{2\sigma^2} \{ \left| \sum_{j=1}^q (y_1^j \, \overline{h}_{2,j} + y_1^j) - y_1^j \right| = const - \frac{1}{2\sigma^2} \{ \left| \sum_{j=1}^q (y_1^j \, \overline{h}_{2,j} + y_1^j) - y_1^j \right| = const - \frac{1}{2\sigma^2} \{ \left| \sum_{j=1}^q (y_1^j \, \overline{h}_{2,j} + y_1^j) - y_1^j \right| = const - \frac{1}{2\sigma^2} \{ \left| \sum_{j=1}^q (y_1^j \, \overline{$$

$$\ln P(x_2 | y_1, y_2, ..., y_q) = const - \frac{1}{2\sigma^2} \Big\{ \Big\| \sum_{j=1}^q (y_1' h_{2,j} + y_2 / h_1' - x_2 + j = 1 q_i = 12 h_i / 2 - 1 x_2 + \ln P(x_2) \Big\}$$

(6)

The corresponding a posteriori probabilities of the bits (i.e. the corresponding soft outputs) using the symbol-to-bit probability conversion of

$$P(d_i = 0) = \sum_j P(x_j | y_1, y_2, \dots, y_q), \quad \forall x_i$$
$$= (d_1 \dots d_i \dots d_b), \quad d_i = 0,$$
(7)

$$P(d_i = 1) = \sum_j P(x_j | y_1, y_2, ..., y_q), \quad \forall x_i = (d_1 ... d_i ... d_b), \quad d_i = 1,$$

Where P(di = 0) or P(di = 1) represents the probability of the ith bit, namely di, of the b-bit symbol being zero and one, respectively. Then the relevant soft outputs can be forwarded to the channel decoders, which will make a hard decision to finally decode the received signals.

3.2 Channel-Coded STBC

The MAP algorithm invoked for decoding STBCs can be exploited by concatenated channel decoders for further improving the system's performance. This paper concatenates on the STBCs with a Turbo Convolutional (TC) code [22, 4]. The STBCs can also be concatenated with a range of other channel codes, such as Convolutional Codes (CCs), Turbo Bose–Chaudhuri–Hocquenghem (TBCH) codes (a class of FEC codes), etc. The best scheme found was the half-rate TC(2,1,4) code in conjunction with the STBC G2. TTCM [23] is a more recent joint coding and modulation scheme which

has a structure similar to that of the family of binary turbo codes,

4. Non-Binary TTCM

Extending binary turbo codes to non-binary turbo codes can be considered, the principle of the non-binary turbo decoding algorithm remains the same. One of the main differences is the trellis diagram associated with a non-binary convolutional code, which has more branches leaving and entering nodes in the trellis, resulting in more paths and higher decoding complexity.

The non-binary turbo encoder has the same structure as the binary turbo encoder, with the component encoders being replaced by RSC codes defined over a ring of integers \mathbb{Z}_M , where M is the cardinality of the ring. The non-binary turbo encoder is given in Figure (3). The message symbols xt and the turbo encoder output symbols $V_t^{(0)}, V_t^{(1)}, V_t^{(2)}$ are now defined over \mathbb{Z}_M as follows:

- $V_t^{(0)} \in \{0, 1, 2, \dots, M-1\}$ is a turbo encoder output corresponding to an information bit at time t.
- $V_t^{(1)} \in \{0, 1, 2, \dots, M-1\}$ is a parity bit from the first component RSC encoder.
- $V_t^{(2)} \in \{0, 1, 2, \dots, M-1\}$ is a parity bit from the second component RSC encoder.

The upper encoder receives the data directly, while the lower encoder receives the data after it has been interleaved by a permutation function, Π . The interleaver is in general a pseudo random interleaver that is it rearranges bits according to a prescribed but randomly-generated rule. The interleaver operates in a block-wise fashion, interleaving L bits at a time, and thus turbo codes can be viewed as block codes.



Fig3: The non-binary turbo encoder

A \mathbb{Z}_M -ring-TTCM encoder system is similar to a classical \mathbb{Z}_M -ring- Turbo encoder (parallel concatenation is evident), but the difference is that, blocks of n coded bits are treated as input symbols, and thus, the interleaver is symbol–oriented, and the component M-ary RSC encoders are trellis encoders–for example, 8–PSK encoder for n = 3, and QAM encoder for n = 4. The final sequence of transmitted symbols is generated by selecting symbols alternately from the two encoders.

4.1 Non-Binary Iterative Turbo Decoding

The idea of the non-binary turbo decoding process is the same idea of the binary turbo decoding process, in which the extraction of extrinsic information from the output of one decoder and pass it on to the second decoder in order to improve the reliability of the second decoder's output and vice versa, as shown in Figure (4).



Fig 4: Iterative decoding in MAP algorithm

The design procedure of a \mathbb{Z}_M -ring-Turbo decoder can be divided into three stages:

- ✓ The first stage is to derive the reliability values of the systematic information, $r_t^{(0)}$, the parity bits from encoder 1, $r_t^{(1)}$, and the interleaved parity bits from encoder 2, $r_t^{(2)}$.
- ✓ The second stage is to employ the Maximum A Posteriori (MAP) decoding algorithm to perform the symbol-by-symbol, which defined over a ring of integersZ_M, decoding and then decision making.
- ✓ The third stage is to derive the mathematical relations that achieve connection between decoder 1 and decoder 2 through iteration cycles.

In non-binary systems, expanding to a ring of integers \mathbb{Z}_M , it must be considered the reliabilities of the other symbols too. The multi-dimensional log-likelihood ratios (multi-dimensional LLRs) for an event u being an element in \mathbb{Z}_M are:

These multi-dimensional LLRs are used by non-binary turbo decoder as its inputs, and their values depend on the type of the channel and the modulation scheme used.

To derive the multi-dimensional LLRs of a 4-state \mathbb{Z}_M -ring-Turbo decoder, with $M \in \{0, 1, 2, 3\}$, and assuming for simplicity, the AWGN channel and 4-ary PAM or 4-ary ASK modulation schemes with constellation points at $(\pm \sqrt{E_s/5}, \pm 3\sqrt{E_s/5})$ are used, where Es, is symbol energy.

Since, the values of the non-binary turbo encoder output symbols are $V_t^{(0)}, V_t^{(1)}, V_t^{(2)} \in \{0, 1, 2, 3\}$, and then mapping of $V_t^{(i)}$ to the 4-ary ASK constellation, $a_t^{(i)}$, i = 0, 1, 2 is given below:

$$a_t^{(0)} = 2V_t^{(0)} - 3, \text{ where } V_t^{(0)} = 0 \text{ or } 1$$

$$a_t^{(1)} = 4V_t^{(1)} - 5, \text{ where } V_t^{(1)} = 1 \text{ or } 2$$

$$a_t^{(2)} = -2V_t^{(2)} + 7, \text{ where } V_t^{(2)} = 2 \text{ or } 3$$

(9)

Thus, $a_t^{(i)} \in \{-3, -1, 3, -1\}$, where i = 0, 1, 2 according to the mapping relations of (4.18).

To calculate the reliability of the systematic information bit, $r_t^{(0)}$:

$$L^{(1)}\left(r_t^{(0)} \middle| a_t^{(0)}\right) = ln \left[\frac{p(r_t^{(0)} \middle| a_t^{(0)} = -1)}{p(r_t^{(0)} \middle| a_t^{(0)} = -3)}\right]$$
(10)

Since, $p(r_t^{(0)}|a_t^{(0)})$ represents the conditional probability density function (PDF) for AWGN channel and is given by

$$p(r_t|a_t) = \frac{(1/\sigma\sqrt{2\pi}) (e^{-(r_t - a_t)^2}/2\sigma^2)}{(1/\sigma\sqrt{2\pi}) (\sum_{i=\pm 1,\pm 3} (e^{-(r_t - ia_t)^2}/2\sigma^2))}$$
(11)

Where σ^2 , represents the noise variance, for 4-PAM modulation with constellation points at $(\pm \sqrt{E_s/5}, \pm 3\sqrt{E_s/5})$:

$$p\left(r_{t}^{(0)} \middle| a_{t}^{(0)} = -1\right) = \frac{p\left(r_{t}^{(0)} \middle| a_{t}^{(0)} = -1\right)}{\frac{(1/\sigma\sqrt{2\pi})\left(\sum_{i=\pm 1,\pm 3}\left(e^{-(r_{t}^{(0)}+i\sqrt{E_{S}/5})^{2}/2\sigma^{2}\right)\right)}{(1/\sigma\sqrt{2\pi})\left(\sum_{i=\pm 1,\pm 3}\left(e^{-(r_{t}^{(0)}+i\sqrt{E_{S}/5})^{2}/2\sigma^{2}\right)\right)}\right)}, \text{ and } p\left(r_{t}^{(0)} \middle| a_{t}^{(0)} = -3\right) = \frac{(1/\sigma\sqrt{2\pi})\left(e^{-(r_{t}^{(0)}+3\sqrt{E_{S}/5})^{2}/2\sigma^{2}\right)}{(1/\sigma\sqrt{2\pi})\left(\sum_{i=\pm 1,\pm 3}\left(e^{-(r_{t}^{(0)}+i\sqrt{E_{S}/5})^{2}/2\sigma^{2}\right)}\right)}\right), \text{ then } L^{(1)}\left(r_{t}^{(0)} \middle| a_{t}^{(0)}\right) = ln\left[e^{-\frac{\left(r_{t}^{(0)}+\sqrt{E_{S}/5}\right)^{2}+\left(r_{t}^{(0)}+3\sqrt{E_{S}/5}\right)^{2}}{2\sigma^{2}}}\right], t 2\sigma^{2} = N.$$

Let
$$2\sigma^2 = N_o$$

 $L^{(1)}\left(r_t^{(0)} \middle| a_t^{(0)}\right) =$
 $(4/5)\left(\sqrt{E_s}/N_o\right)r_t^{(0)} + (8/5)(E_s/N_o)$ (12)

Thus, each one of the systematic information bit, $r_t^{(0)}$, the parity bit from encoder $1, r_t^{(1)}$, and the interleaved parity bit from encoder $2, r_t^{(2)}$, has three reliability values, respectively, as shown below in system of equations:

$$L^{(i)}\left(r_{t}^{(0)}|a_{t}^{(0)}\right) = (4/5)\left(\sqrt{E_{s}}/N_{o}\right)r_{t}^{(0)} + (8/5)(E_{s}/N_{o}),$$

$$L^{(i)}\left(r_{t}^{(1)}|a_{t}^{(1)}\right) = (12/\sqrt{5})\left(\sqrt{E_{s}}/N_{o}\right)r_{t}^{(1)}, \quad (13)$$

$$(i)\left(r_{t}^{(2)}|a_{t}^{(2)}\right)$$

$$L^{(0)}(r_t^{(n)}|a_t^{(n)}) = (8/\sqrt{5})(\sqrt{E_s}/N_o)r_t^{(2)} + (8/5)(E_s/N_o).$$

Where i = 1, 2, 3.

Thus the output of each decoder, in each decision case is defined by the multi-dimensional LLR: Case 1: the decision between (0&1):

$$L^{(1)}\left(\left(a_t^{(0)}\middle|r\right) = ln\left(\frac{P(a_t^{(0)} = -1|r)}{P(a_t^{(0)} = -3|r)}\right)$$
(14)

Case 2: the decision between (0&2); $L^{(2)}((a_t^{(0)}|r)) = ln\left(\frac{P(a_t^{(0)}=+3|r)}{P(a_t^{(0)}=+3|r)}\right)$

$$L^{(2)}\left(\left(a_{t}^{(0)} \middle| r\right) = ln\left(\frac{1}{P(a_{t}^{(0)} - 3|r)}\right)$$
(15)

$$L^{(3)}\left(\left(a_{t}^{(0)}\middle|r\right) = ln\left(\frac{P(a_{t}^{(0)}=+1\middle|r)}{P(a_{t}^{(0)}=-3\middle|r)}\right)$$

(16)

5. Proposed Z_M-ring-TTCM Modulationassisted STBC-OFDM Scheme

 \mathbb{Z}_4 -Ring-TTCM channel code can be combined with the G2 STBC for the sake of improving the system's performance. Furthermore, the above \mathbb{Z}_4 -ring-TTCM-assisted G2 STBC-aided scheme would be concatenated with a single-user OFDM system. Figure (5) shows a schematic of the \mathbb{Z}_4 -ring-TTCM-assisted G2 STBC-OFDM system. The source information bits are first encoded and modulated by the \mathbb{Z}_4 -ring-TTCM encoder followed by the space-time encoder, the STBC employed was the G2 code, which invokes two transmitter antennas, and the two space-time-coded samples are mapped to two consecutive OFDM subcarriers and OFDM modulated. The OFDM symbols are then transmitted via the multi-path fading channel, and the received noise-contaminated symbols are forwarded to the OFDM demodulator. The recovered signal is then space-time soft decoded and the soft outputs are fed to the \mathbb{Z}_4 -ring-TTCM-decoder which used the multi-dimensional MAP algorithm for recovering the most likely transmitted information bits.



Fig 5: Schematic diagram of the proposed \mathbb{Z}_4 -ring-TTCMassisted G_2 STBC-OFDM system.

6. Simulation Results

The schematic design of \mathbb{Z}_4 -ring-TTCM encoder can be simulated on computer. The \mathbb{Z}_4 -ring-TTCM encoder passes through 4-state transitions, since at each state two sets of nonbinary mathematical operations are used to encode a message, one for the normal turbo code output and the other for the interleaved output. A modulation scheme, pulse amplitude modulation of fourth order, is suitable to modulate the encoded bits with low probability of errors. The simulation results of both binary TTCM-QPSK-assisted G2 STBC-OFDM [12] and the \mathbb{Z}_4 -Ring-TTCM scheme-assisted G2 STBC-OFDM can be shown in Figure (6), where the code rate of G2 STBC scheme is (1) and code rate of TTCM is (1/2).



Fig 6: The performance comparison of (BTTCM and Z_4-Ring-TTCM)-assisted G2 STBC-OFDM schemes

The performances of the uncoded BPSK scheme [13], the G2 STBC scheme [24], the BTTCM-QPSK scheme-assisted G2 STBC-OFDM system [12], and the \mathbb{Z}_4 -Ring-TTCM-PAM scheme-assisted G2 STBC-OFDM system, can be summarized in Table (1), where the coding gains are defined

as the (Eb/No) difference, expressed in decibels, at BERs of 10^{-5} and 10^{-3} between the various channel coding assisted STBC-OFDM systems and the uncoded single transmitter system having the same effective throughput. The performance of the best scheme in Table (1) is (printed in bold), since the performance comparison shows that the Z₄-ring-TTCM-PAM scheme-assisted G2 STBC-OFDM outperforms the systems [13, 12, 24]. Since, the Z₄-ring-TTCM-PAM scheme-assisted G2 STBC-OFDM system is provided that the gains in (Eb/No) are (18.60dB) and (35.12 dB) at the BERs of 10^{-3} and 10^{-5} respectively as shown in Table (1). The effective throughput of the Z_P-ring-TCM-PAM and Z_P-ring-TTCM-PAM can be expressed by:

$$\eta = R\left(\sqrt{\mathbf{p}} \cdot \log_2 M\right) \tag{17}$$

Where R, is the code rate, $p \in \mathbb{Z}$, and M, is the modulation order. Since, for R = 1/2, p = 4, and M = 4, then η = (2) bits/s/Hz.

7. Conclusions and Future Works

It is shown that the idea TTCM codes has been extended for symbols defined over rings of integers \mathbb{Z}_q and that increased the bandwidth efficiency by a factor of \sqrt{q} over the binary TTCM codes. The use of non-binary TTCM codes led to reduction in the effective input block length, since each *m* bits of binary information correspond to one non-binary symbol for $q = 2^m$, and thus non-binary system can be used with high number of symbols.

Table (1) the performance comparison between binary and non-binary CM schemes-assisted G2 STBC-OFDM system.

		BER				Rayleig	
			$E_b/N_o(dB)$		Gain (dB)		h fading
STBC scheme	CM scheme	CM Code rate	10-3	10-5	10-3	10 ⁻⁵	Modem
Uncode d			24.1 1	44.1 2	0.00	0.00	BPSK
G_2			13.9 3	25.8 7	10.1 8	18.2 5	BPSK
G2 STBC- OFDM	TTCM	1/2	7.91	11.8 2	16.2 0	32.3 0	QPSK
G ₂ STBC- OFDM	Z ₄ - RTTC M	1/2	5.51	9.00	18.6 0	35.1 2	4-PAM

Non-binary TTCM scheme that has modulation order (M) can achieve an error performance similar to that of binary schemes that have higher order (M), and this is the reason of achieving good performance by non-binary systems over binary systems. Non-binary turbo decoding was achieved by introducing an array of LLR values for each non-zero element in the ring, instead of just one in binary decoding. The drawbacks of non-binary Turbo codes are; more branching leaving each state of trellis structure, non-binary symbols and LLR values, and more computations complexity that need more storage memory. It is shown that Ring-TTCM codes are suitable for combining with MIMO-OFDM system, the performance of these codes were evaluated well than the binary codes, and hence, these codes achieve the 4G wireless candidate of efficient bandwidth and good error control performance.

8. REFERENCES

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