Design and Implementation of an Efficient Instruction Set for Ternary Processor

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ABSTRACT

Multi Valued Logic [MVL] is emerging as a promising choice for future computing technology. MVL has seen major advancement in the recent past due to several advantages offered by them over the binary logic, thus making it a thrust area for further research. The instruction set of the processor is its inherent entity. This paper presents design and implementation of an efficient instruction set for a ternary processor using Very-High-Speed Integrated Circuits, VHSIC Hardware Description Language [VHDL]. Twenty one instructions including various addressing modes such as register, direct and immediate mode are designed and implemented for 4-trit ternary processor. The required control signals are appropriately identified in the proposed design and enable the smooth operation of instructions. The designed 4 trit instruction set signifies encouraging results that will pave the path for further developments in ternary processors.

General Terms

Digital Systems

Keywords

Multi Valued Logic, Ternary logic, VHDL.

1. INTRODUCTION

Multi Valued Logic [MVL] has revolutionized the future of digital systems. The switching algebra in digital systems has expanded from binary to penta level, progressing through ternary and quaternary levels. *Elena et al.* describes MVL as painting a picture having all possible colors available, as against binary logic which includes just black and white. MVL displays a phenomenon that is never seen in binary, where the only two possible values are null and unity elements of Boolean algebra, possessing very specific properties [1]. An excellent review of the first developments of MVL and its applications to electronics, including a few circuit implementations, was presented by *Epstein et al.* [2]

In a VLSI circuit approximately 70 % of the area is devoted to interconnection, 20 % to insulation and 10 % to devices. The binary logic is limited due to interconnect which occupies large area on VLSI chip [3]. Ternary Logic offered a potential solution to one of these major problems in VLSI System [4]. Advantages of ternary system over binary system are:

- Reduced number of interconnections. For example, $(72)_{10} \Rightarrow (1100)_3 \Rightarrow (1001000)_2$.
- Reduced chip area.
- Reduction in chip area in turn reduces chip delay.
- More information can be transmitted. For example,

Using 4 bits, only 0 to 15 numbers can be represented in binary system while 0 to 80 in ternary system.

- It also offers better utilization of transmission channels because of the higher information content carried by each line.
- It gives more efficient error detection and correction codes.
- Serial and serial parallel arithmetic operations can be carried out faster by using the ternary logic.
- One of the main advantages of ternary logic is that it reduces the number of required computation steps.
- The number of address lines required to access the address content can be reduced, i.e. hardware cost can be reduced.
- Less memory requirement.

As the value of radix increases, the information carrying capacity of each connection also increases. Hence, MVL logical systems, for instance, a three-valued (radix 3) digital realization would be more appropriate than binary. Ternary (or three-valued) means a switching element, which switches among 3 levels namely true, false and intermediate. As ternary is radix-3 system, three levels are considered as logic '0' = 0V, logic 'Z' = 2.5V and logic '1' = 5V. So, this ternary processor should be capable of generating control signals for the instructions to be executed for a 4-trit microprocessor. In short, using ternary logic or radix-3 system, the number of address lines required to access the address content can be reduced, i.e. hardware cost can be reduced and more data can be occupied within less memory space. An evident advantage of a ternary representation over binary is economy of digits. To represent a number in binary system, one needs 58% more digits than that of ternary. It is therefore that the ternary is casting its applications in the field of Fuzzy logic, Machine Learning, Artificial Intelligence, Data Mining, Robotics, Digital signal processing, Digital control systems and Image Processing.

Ternary logic has attracted a wide research community to contribute towards building an efficient ternary processor. *Sheng Lin et al.* proposed a ternary CNTFET based design of ternary logic gates and arithmetic circuits [6], *Staishkumar et al.* [7] modeled combinational circuits based on ternary multiplexers using VHDL. *Milton et al.* [8] developed MVL algebra for the synthesis of digital circuits, whereas *Raychowdhury* [9] proposed a novel multiple valued designs using ballistic carbon-nano-tube FET. Many fundamental building blocks like ternary gates, adders, ALU and memory elements of ternary processor are also designed. *Mozamel et al.* [10,11] devised a quantum ternary parallel adder / subtractor with partially look ahead carry whereas *Dhande et al.* [12] designed and implemented 2 bit ternary ALU slice

using CMOS ternary logic gates. *Ali et al.* [13] exploited the usefulness of neural network for designing a ternary ALU. Furthermore *Kotiyal et al.* [14] extended the reversible logic for a ternary barrel shifter. Use of various EDA tools like VHDL [15] for the simulation of ternary circuits have been also been reported in the literature.

The literature thus reports extensive efforts to develop combinational and sequential circuits using ternary logic. These designed circuits are certainly the essential components of a ternary processor. However, additionally, an efficient instruction set also forms an imperative component and plays a vital role in ternary processor. This issue has received comparatively less attention by the researchers and meager work is reported in the literature addressing this concern, to the best of the knowledge.

This paper proposes a novel approach to design and implement an efficient instruction set for ternary processor using VHDL. Twenty one instructions including various addressing modes such as register, direct and immediate mode are considered during this research. Section 2 describes the design of the proposed instruction set and the next section details the simulation results. Paper concludes with a discussion of the results and further scope for enhancing the designed instruction set.

2. DESIGN OF THE INSTRUCTION SET

The objective of the proposed research is to design a 4-trit instruction set for a ternary processor using VHDL as a simulation tool. VHDL supports a top down and a bottom up approach by allowing higher level system design. VHDL also supports nine states, defining various logic levels like U, X, 0, 1, Z, W, L, H and -, that represent logic levels namely: uninitialized, unknown, logic 0, logic 1, high impedance, weak unknown, weak zero, weak one and don't care, respectively [16]. The user defined packages, libraries and the user defined data types are the significant features of VHDL that are exploited for the design and simulation of the ternary instruction set which involves consideration of multiple logic levels.

VHDL can be used to model, simulate and describe ternary system where signals in the circuit can take tristate logic i.e., the usual true and false states, with a third transparent high impedance state. The third (intermediate) state can be metaphorically thought of as either an unambiguously true or unambiguously false. To demonstrate the use of VHDL as a ternary logic simulator, logic-0 is used to represent 0 volt, high impendence Z to represent 1volt and logic-1 to represent 2 volts. VHDL provides an effective way to connect several logic outputs to a single input, where all but one is forced to the high impedance state, allowing the remaining outputs to operate in the normal binary sense. This concept is commonly used for memory bank connection in computers and in connecting other similar devices to a common data bus. A large number of devices can therefore communicate over the same channel, simply by ensuring that only one is enabled at a time. VHDL package which include functions for ternary logical operations such as T-NOT, T_AND, T_NAND, T_OR, T_NOR, T_XOR are build successfully. It also includes ternary clock with clock period of 40 ns. Logical and arithmetic functions are defined in VHDL package and named as 'std terlogic'. The designed instruction set includes the instructions of various addressing modes namely. Register, Direct and immediate, as indicated in Table 1.

The assigned op-codes and an example of each instruction are also depicted for comprehensiveness. Ternary instruction decoder is also designed that will decode the three trit op code from IR register. Instruction decoder is 3:27 decoder, which will select the corresponding instruction amongst the set of twenty one instructions. The control signals that drive the instruction set are also considered for implementation.

Table 1: The proposed instruction set for ternary processor.										
			-	-	Addressing	_				

Sr. No	Mnemonic	Opcode	Addressing Mode	Example
1.	T_ANA	00Z	Register	T_ANA B
2.	T_ORA	001	Register	T_ORA B
3.	T_XRA	0Z0	Register	T_XORA B
4.	T_ADD	0ZZ	Register	T_ADDA B
5.	T_ADC	0Z1	Register	T_ADC B
6.	T_SUB	010	Register	T_SUB B
7.	T_SBB	01Z	Register	T_SBB B
8.	T_ICR	011	Register	T_ICR
9.	T_DCR	Z00	Register	T_DCR
10.	T_RAL	Z0Z	Register	T_RAL
11.	T_RAR	Z01	Register	T_RAR
12.	T_CMA	ZZ0	Register	T_CMA
13.	T_ADDI	ZZZ	Immediate	T_ADDI, 35
14.	T_SUI	ZZ1	Immediate	T_SUI, 45
15.	T_MVI A	Z10	Immediate	T_MVI A,20
16.	T_MVI B	Z1Z	Immediate	T_MVI B, 20
17.	T_LDA	Z11	Direct	T_LDA, 20
18.	T_MOV A, B	100	Register	T_MOV A,B
19.	T_MOV B, A	10Z	Register	T_MOV B,A
20.	T_MOV A, C	101	Register	T_MOV A,C
21.	T_MOV C, A	1Z0	Register	T_MOV C,A



Figure 1: Data flow diagram of the ternary processor.

3. DATA FLOW AND SIMULATION RESULTS

Figure 1 depicts the dataflow diagram for the proposed instruction set. Ten different control signals are generated as explained below:

PC_sig	:	To control in and out operation of the PC register
Y_Z_sig	:	To control in operation of Y register and out operation of Z register
R_W_sig	:	To control read and write operation from/on memory
A_sig	:	To control in and out operation of the A register
B_sig	:	To control in and out operation of the B register
C_sig	:	To control in and out operation of the B
MAR_sig	:	To control in and out operation of the MAR
MDR_sig	:	register To control in and out operation of the MDR register
IR_sig	:	To control in and out operation of the IR register
END	:	To specify the end of the instruction execution
selmux	÷	To select either ALU or a constant
sel(0)	:	First trit (LST) of ALU select
sel(1)	:	Second trit of ALU select
sel(2)	:	Third trit (MST) of ALU select

Few control signals will control the register dataflow i.e in and out. For such a trit control signal,

When signal is '1' \Rightarrow register out operation When signal is 'Z' \Rightarrow register in operation When signal is '0' \Rightarrow No operation

A control signal specifies the read and write operation for memory.

When signal is '1' \Rightarrow memory read operation When signal is 'Z' \Rightarrow memory write operation When signal is '0' \Rightarrow No operation

Three trit control signals are used to select ALU operation i.e. sel(0), sel(1), sel(3). Selmux is the control signal which will select the one of the input to the ternary ALU. Input can be either accumulator or a constant 'Z'. END signal will specify the end of the instruction and it will reset the two trit step counter. Processor can execute up to twenty three instructions. These instructions are performed in basically three steps:

- Fetch: Fetch the ternary op-code from the memory and send it to the instruction register for decoding.
- Decode: Decode the op-code in the instruction register and decide which instruction is to be executed.
- Execute: Execute the decoded instruction by activating the appropriate control signals' sequence.

The dataflow diagram of the processor enables the working of the processor. Consider the instruction " $T_MVI A$, 00Z1" which is an instruction with immediate addressing mode. The opcode of the instruction "0Z10" is stored in the memory at one location and the immediate data "00Z1" is stored at the next location. In this instruction the immediate data "00Z1" is copied from the memory to the accumulator. The instruction is performed in a total of seven steps as shown below:

• Step 1:

Control signals: PC=>'1', MAR=>'Z', Y_Z=>'Z', sel_mux=>'1', sel_ALU=>"0Z1", R_W=>'1'. Operations performed:

First, the address of the opcode is transferred from the PC to MAR register and read signal is activated. The data then is read into the MDR register in the same cycle. At the same time, value in program counter is incremented in the ALU by "000Z".

• Step 2:

Control signals: Y_Z=>'1', PC=>'Z'. Operations performed: The incremented value of program counter is updated.

Step 3:

Control signals: MDR=>'1', IR=>'Z'. Operations performed: The opcode fetched into the memory data register is now copied into the instruction register.

Step 4:

Control signals: PC=>'1', MAR=>'Z', Y_Z=>'Z', sel_mux=>'1', sel_ALU=>"0Z1", R_W=>'1'. Operations performed: Similar to step 1, the data at the next location, that is, the immediate data is now fetched and the value in program

counter is again incremented in the ALU.

Step 5:

Control signals: Y_Z=>'1', PC=>'Z'. Operations performed: The incremented value of program counter is again updated.

Step 6:

Control signals: MDR=>'1', A=>'Z'. Operations performed: The fetched immediate data is copied from the memory data register into the accumulator.

• Step 7:

Control signals: End=>'1'.

Operations performed:

This indicates that the instruction has completed execution. Figure 2 (a)-(f) illustantes the step wise results indicating the described dataflow. Figure 3 depicts the simulation results of this instruction.



Figure 2(a): Step 1, T1 cycle



Figure 2(b): Step 2, T2 cycle



Figure 2(c): Step 3, T3 cycle



Figure 2(d): Step 4, T4 cycle



Figure 2(e): Step 5, T5 cycle



Figure 2(f): Step 6, Cycle T6

4. CONCLUSION

This paper proposes an efficient design of 4 trit instruction set for twenty one instructions with various addressing modes such as register, direct and immediate mode. The required control signals are appropriately identified in the proposed design and enable the smooth operation of instructions. The designed 4 – trit instruction set signifies encouraging results that will pave the path for further developments in ternary processors.

The issue of effectively designing the instruction set for a ternary processor has received comparatively less attention by the researchers and meager work is reported in the literature addressing this concern. It must therefore be emphasized that the performance of the proposed ternary instruction set though is extremely encouraging; it cannot be directly compared with the state of art methods. The depicted simulation results provide a strong foundation about the behavior of the ternary processor from the instruction – execution prospective; which is governed by the identified control signals.

Our further research will focus on enhancing the proposed design to incorporating a greater number of instructions like multiplication and division operations and increasing the number of addressing modes like indirect addressing mode, index addressing mode. The word size will also be expanded to more than 4-trit and the range of data that can be processed will be increased. Further improvement may also include different types of addresses involve tradeoffs between instruction length, addressing modes, flexibility and complexity of address calculation.

— ternary Clock ——	0																				
- Instruction register -	·		- I																		
∎-4) ir_bus	071	071	\exists	Z10							102					0ZZ					
- registers	_																				
₽-♦ pc_bus	1200	1200	1202			1201				1220					1772						177
∎-<> mar_bus	1200	1200			1202				1201					1220						1772	
∎-<> mdr_bus	0Z10	0710			0011				0102					0022						0100	
∎-♦ y_bus	1200	1200			1202				1201					1220			0011			1772	
0-4 z.bus	1202	1202	7		1201				1220					1222			0712			1221	
∎-🔶 a_bus	0000						0011											0Z1Z			
∎-今 b_bus	0011	0011										0011									
🖬 🕁 c.bus	0000																				
— Control signals —																					
📣 pc	1		~	-																	
📣 ir	0																				
4 y_2	Z																				
👌 run	1		٦.																		
4 mar	Z																				_
🔥 mdr	0																				
👍 selmux	1																				
🔥 sel_2	0																				
👌 sel_1	z									_											
sel 0	1					-7															7
4 a	0							_											_		
l do b	0																				
Å.	0																				
Now 101000000 p	ы К	Ľ	<u>ا</u>	990000	00 ps	1.1	1.1	99200	000 ps	1.1	1.1	99400	000 ps	1.1	1.1	99600	000 ps	1.1	1 1	1 1 1 998000	00 ps

Figure 3: Simulation results of the instruction execution

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