Achieving *k*-anonymity using Minimum Spanning Tree based Partitioning

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ABSTRACT

Protecting individual's privacy has become a major concern among privacy research community. Many frameworks and privacy principles were proposed for protecting the privacy of the data that is being released to the public for mining purpose. k-anonymization was the most popular among the proposed techniques in which the sensitive association between the sensitive attributes and their corresponding identifiers are de-associated. In this paper, we proposed an enhanced k-anonymity technique by using Minimum Spanning Tree (MST) partitioning approach. In this technique we disclose the information of the individuals pertaining to minimum group size i.e., k. We achieve this technique in two phases. During the first phase, MST for the given dataset is partitioned to generate equivalence classes and in the subsequent phase whether the equivalence class size is achieved to that of the minimum group size k is verified. Our approach resulted in achieving the optimal anonymization along with data utility. We showed the efficacy of our proposed technique by running a series of experiments in terms of information loss to show that our technique adheres to the quality of the anonymized data.

General Terms

Privacy Preserving Data Publishing, Anonymization, Algorithm.

Keywords

Privacy, Anonymisation, Hierarchical distance, Minimum Spanning Tree, Inflexion point.

1. INTRODUCTION

Several data holders such as hospitals, financial corporations and other organizations publish their microdata for the purpose of data mining, statistical analysis, and other public benefits. Publishing data as it is may reveal the person specific data. The attributes of the data table are divided into three categories, namely identifying attributes, quasiidentifiers and sensitive attributes. The attributes like "Name", "SSN" etc., from which we can identify the individuals directly are termed as identifying attributes. The attributes like "Age", "Gender" etc., from which we can potentially identify the individuals are termed as Quasi Identifiers. The attributes like "Salary", "Disease" are defined as sensitive attributes. perturbation, anonymity frameworks like generalization, suppression and cryptographic techniques are being used for protecting the microdata [15, 16, 1, 2, 3]. In perturbation, data is de-identified by adding noise to original values before publishing. However, the values of the

perturbed data values are slightly deviated from the original data values. Generalization replaces more specific value with

a less specific value [3]. In cryptographic techniques the data is encrypted such that any party cannot see other parties' data during data sharing [15, 16].

The traditional approach is to de-identify the microdata by removing identifying attributes such as name, SSN etc. [1]. Sweeney discovered that when the medical data was linked to the external voters list, 87% of the US citizens were identified [3]. This type of attack is coined as linking attack. To limit these kinds of attacks Sweeney proposed k-anonymity model where the domain of each quasi identifier attribute is divided into set of equivalence classes and each equivalence class contains at least k elements with the same value [2, 3]. In this model the quasi identifier attributes are suppressed or generalized. The data is generalized by constructing the Domain Generalization Hierarchies (DGH) for the corresponding quasi identifiers. For instance the DGH for age, gender and zip-code are shown in Fig 1(a), 1(b), 1(c) respectively. For the original Table 1(a), a 2-anonymised table is shown in Table 1(b) and 3-anonymized view is shown in Table 1(c). Here sensitive attributes are retained without any change.

Table 1(a). Original Table

ID	Age	Sex	Zip-code	Disease
t_1	21	M	13051	Anaemia
t_2	26	M	13053	Flu
t_3	33	F	13063	Cancer
t ₄	37	F	13068	HIV
t ₅	45	M	14851	Flu
t ₆	48	F	14856	HIV
t ₇	50	M	14865	Cancer

Table 1(b). 2- Anonymized table

ID	Age	Sex	Zip-code	Disease
t_1	[20-30]	M	1305*	Anaemia
t_2	[20-30]	M	1305*	Flu
t_3	[31-40]	F	1306*	Cancer
t ₄	[31-40]	F	1306*	HIV

t ₅	[41-50]	Person	148**	Flu
t ₆	[41-50]	Person	148**	HIV
t ₇	[41-50]	Person	148**	Cancer

Table 1(c). 3- Anonymized table

ID	Age	Sex	Zip-code	Disease
t_1	[20-40]	Person	130**	Anaemia
t_2	[20-40]	Person	130**	Flu
t ₃	[20-40]	Person	130**	Cancer
t ₄	[20-40]	Person	130**	HIV
t ₅	[41-50]	Person	148**	Flu
t ₆	[41-50]	Person	148**	HIV
t ₇	[41-50]	Person	148**	Cancer

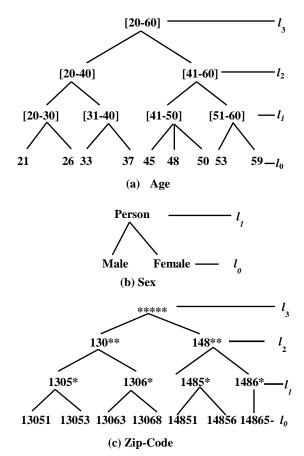


Fig 1: Domain generalization hierarchies

In this paper we proposed minimum spanning tree based technique to achieve *k-anonymity*. In this approach the minimum spanning tree is constructed by considering quasi identifiers as the vertices of the spanning tree and the edges of the spanning tree represents hierarchical distance between two tuples. Initially, MST forest is constructed for the given microdata. We then remove all inconsistent edges which result in forming sub-trees. This process is repeated until each

subtree has *k* data points. Once this is achieved we generalize the data points adhering to *k-anonymity* for publishing.

The rest of the paper is organised as follows: Section 2 surveys the related work. Section 3 presents an overview of minimum spanning tree. In Section 4 we present some formal definitions that were used in this paper. MST based anonymisation model in explained in section 5. Section 6 presents algorithms and discusses the complexity measures of our approach. The essential quality measures necessary for assessing our method are given in section 7. Section 8 show experimental evaluations to show the quality of our apporach and finally we conclude and present, possible future work in section 9.

2. RELATED WORK

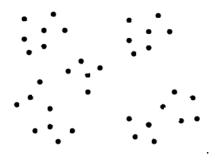
Sweeney and Samarati introduced k-anonymity, in which the domain of each quasi identifier attribute is partitioned into set of intervals by replacing the attribute values with corresponding intervals [2, 17]. Here, the tuple in the published table is indistinguishable from at least k-1 other tuples with respect to their quasi identifiers. Various models like global recoding, local recoding, multidimensional recoding, micro aggregation and clustering were proposed to achieve k-anonymity principle [14, 12, 6, 11]. In global recoding, all the attribute values in generalized table are obtained from the same level of domain hierarchy [20]. The advantage of the global recoding is that anonymized data can be viewed uniformly [19]. However, the anonymized dataset suffers from more information loss. In local recoding, the attribute values of generalized table are generalized to different levels in domain hierarchy. Multi-dimensional kanonymity model assumes a d-dimensional domain space, which is divided into set of regions where each region contains at least k objects.

Another approach to accomplish k-anonymity is micro aggregation [18, 11]. Micro aggregation can be viewed as clustering problem where the size of the cluster is a constraint. Hua Zhuand Xiaojun showed k-anonymity as density based clustering problem in which set of tuples are grouped based on k-nearest-neighboring distance [5]. GaganAgarwal et.al proposed a method called anonymity via clustering [6]. In this method, initially quasi identifier records are clustered and then cluster centers are published. They considered r-GATHER as a metric to form clusters. A frame work called KACA to accomplish the k-anonymity in which grouping of the tuples is done based on attribute hierarchical structures [4]. A non-homogenous generalization for partitioning the data was also proposed that adheres to k-anonymity principle [19]. They enhanced this algorithm by proposing a randomized algorithm such that it does not compromise the attacker even if the algorithm is known [23].

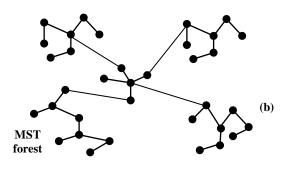
A practically feasible algorithm called enhanced *k-anonymity* was proposed [12]. Their algorithm is based on full domain generalization approach. Fung et al. proposed classification model to protect the individual privacy [13]. They use top-down approach to achieve the anonymity constraint. Their model supports both numerical and categorical data. Graham et al. proposed bipartite graph model to anonymized the data using (*k*, *l*) grouping strategy thereby minimizing the adversary attacks [25,9].

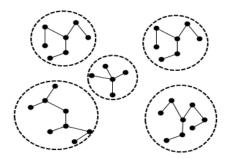
3. MINIMUM SPANNING TREE PARTITIONING

Given a connected, weighted, undirected graph G = (V, E) together with a weight function $c: E \to \Re$, where V is the set of vertices and E is the set of edges between the pair of vertices and c (e) specifying the weight of each edge. A spanning tree of G is a sub-graph which is a tree that connects all the vertices in G without having any cycles. The weight of a spanning tree is the sum of the weights of edges in that spanning tree. A minimum spanning tree (MST) of G is a spanning tree whose weight is less than or equal to weight of every other spanning tree of G. Well known algorithms for finding MST are Kruskal's algorithm [21], Borůvka's algorithm [26], and Prim's algorithm [22].



(a) Data points





(c) Sub-trees of MST forest

Fig 2: MST Partition model

For example, in Fig 2(a) sample data points are considered as the nodes of the graph and distances between these data points are edges of graph. By using the Kruskal algorithm we construct MST forest from the graph. Fig 2(b) show the MST forest for the given data points. After MST is constructed we remove the inconsistent edges from MST. Inconsistent edges are defined as those edges whose edge distances are significantly larger than average distance of the nearby edges

of MST. Hence the MST forest is partitioned into set of subtrees as shown in Fig 2(c). The data points in each partitioned sub-tree are close to each other.

4. PRELIMINARIES

The basic concepts and terminology used in the paper are discussed in this section. Let T be the microdata table that needs to be published. The table contains collection of tuples from domain $D = I \times QI \times SA$, where I be the identifying attributes, QI be the quasi identifiers and SA be the sensitive attributes. For each tuple t ϵ T, the attribute value is denoted as t [A].

Definition 1 (**Generalization**) Let A_i be QID attribute and D_{A_i} be the domain of A_i . Suppose $d_{A_i} \in D_{A_i}$ is an instance of the attribute A_i is said to be generalized to $d_{A_i}^{(l)}$ if and only if there exists a binary relation $\leq_{A_i}^{f_i^{(l)}}$ such that $d_{A_i}^{(0)} \leq_{A_i}^{f_i^{(l)}} d_{A_i}^{(l)}$, where $f^{(l)}$ is I^{th} level of generalization over hierarchy tree and $d_{A_i}^{(l)}$ is I^{th} level generalized value of d_{A_i} . i.e

$$D_{A_i}^{(l)} = \{d_{A_i}^{(l)} | d_{A_i}^{(k)} \leq^{f_i^{(l)}}_{A_i} d_{A_i}^{(l)}, \forall d_{A_i} \in D_{A_i}, k = 0, 1, 2 \dots l - 1\}$$

Definition 2 (*k*-anonymity) The table is said to be *k*-anonymous if every combination of quasi identifier attribute values in a generalized table occurs *k* or more times.

5. FRAMEWORK FOR MSTK

The proposed framework minimum spanning tree partitioning (MSTK) is shown in Fig 3. This framework consists the following phases:

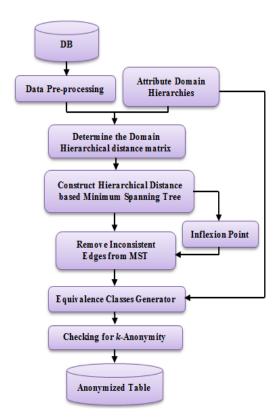


Fig 3: MSTK Framework

Step 1: Data Pre-processing and corresponding domain generalization hierarchies are created.

Step 2: Computing the hierarchical distances among data

Points of given micro table to Generate distance matrix.

- Step 3: Constructing the minimum spanning tree based on the distance matrix.
- **Step 4:** Remove inconsistent edges from MST and form basic equivalence classes.
- Step 5: Checking for anonymity of each equivalence class.

5.1 Data pre-processing

In this step the quasi identifiers attributes which are to be anonymized and their concept hierarchies are selected. For example Sex, Age and Zip-code attributes are quasi identifier attributes in the Table. 1 (a) and their domain hierarchies are shown in Fig 1.

5.2 Computing hierarchical distance matrix

In this section we compute the distance matrix among all the tuple by using the following definitions

Definition 3 (Generalization Hierarchical Distance (GHD)): Let v and v^I be the two nodes of attribute generalization hierarchy and H be the height of the tree. The generalization hierarchical distance between the two node are defined as

$$GHD(v,v') = \frac{|v-v'|}{H} \tag{1}$$

For example the values of an attribute Zip-code in hierarchy are {13051, 1305*, 130**, *****} as shown in Fig.1(c). Distance between 13051 and 130** is 2/3. The GHD is zero if both values are at same level or at the same leaf nodes and GHD is one if the value lies at the root.

Definition 4 (Generalization Effort (GE)): Let t and t¹ be the tuple and generalized tuple respectively. The generalization effort of a tuple is defined as the amount of effort needed to change the attribute values of tuple one (low) level to another (generalized) level in domain generalization hierarchies of attributes of a tuple. i.e.

$$GE(t,t') = \sum_{i=1}^{m} GHD(a_i,a_i')$$
 (2)

Here, a_{i_1} , a_{i_1} are the original attributes values and generalized attribute values of the tuple t respectively. For example, consider tuple $t_2 = \{26, M, 13053\}$ in the Table 1(a) and its generalized values $t_2' = \{[20\text{-}30], M, 1305*\}$ in Table 1(b). The GHD (26, [20 - 30]) = 0.333, GHD (M, M) = 0 and GHD (13053, 1350*) = 0.333, therefore $GE(t_2, t_2') = 0.666$.

Definition 5 (Hierarchical Distance between two tuples): Let t_1 and t_2 be the two tuples and t_{12} is the closest common generalization of tuples t_1 and t_2 . The hierarchical distance between two tuples is defined as follows.

$$HDist(t_1, t_2) = GE(t_1, t_{12}) + GE(t_2, t_{12})$$
 (3)

For example, consider two tuples the tuples $t_1 = \{21, Male, 13051\}$ and $t_2 = \{26, Male, 13053\}$ in the Table 3.1 (a) and their closest common generalization attribute values of tuples are $t_{12} = \{[20\text{-}30], M, 1305*\}$, which is obtained from the generalization attribute hierarchies of age, Sex and Zip-code as shown in Fig. 1 (a), (b), (c). Therefore the Hierarchical distance between two tuples is 1.333.

Definition 6 (Dissimilar Distance Matrix of Table): Given microdata T with n tuples $\{t_1, t_2, ..., t_n\}$ and each tuple contain m quasi-identifiers, The distance matrix is defined as follows and the distance matrix shown in Table 2:

$$D_{T} = \left[Hdist(t_{i}, t_{j})\right]_{n \times n} \forall i, j \in n$$
 (4)

Table 2. Distance Matrix

Γ 0	_	_	_	_	_	-1
1.333	0	_	_	_	_	-1
4.667	4.667	0	_	_	_	-1
4.667	4.667	1.333	0	_	_	-
4.000	4.000	6.000	6.000	0	_	-1
6.000	6.000	4.000	4.000	3.333	0	-1
$L_{4.000}$	4.000	6.000	6.000	2.000	4.000	01

5.3 Minimum Spanning Tree Forest Construction

After computing the distance matrix, MST is created by using above distance matrix. The nodes of the MST are the data points (records) of the micro table and the weight of edges are the concept hierarchical distance between two data points. We adopt Kruskal algorithm for constructing the MST. For example, the MST for the Table 1(a) is shown in Fig 4. and the corresponding weights are given in Table 3.

Table 3. Weights of the MST

Edge	1-2	3-4	3-6	5-6	5-7	1-5
Weight	1.33	1.33	4.0	3.333	2.0	4.0

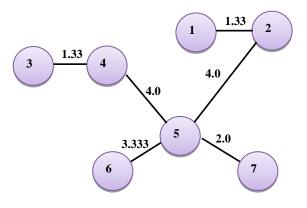


Fig 4: MST forest

5.4 Pruning of inconsistence Edges from MST Forest

For partitioning the MST, the inconsistent (longest) edges are pruned. This pruning process is performed based on inflection point. Finding the inflection point of the data is given as follows:

Inflexion point of the data: Let f be real values function i.e. $f: \mathcal{D} \to \Re$ where \Re is the set of real values and \mathcal{D} is the domain of f. Inflexion points are the points at which the concavity changes. Concave up corresponds to positive second derivative of the curve f, whereas Concave down corresponds to negative second derivative. So, the point at which the second derivative is equal to zero turns the concavity up to down and vice versa. Therefore at inflexion point, the curves' second derivative is equal to zero, i.e., f'' = 0

If we consider the weights of the edges of a spanning tree to be a value of a random variable \mathcal{X} . We assume that the weights of the spanning tree are normally distributed ($\mathcal{X} \sim N(\hat{w},\sigma)$), where \hat{w} be the mean and σ be the standard deviation. If f(x) is probability distribution function \mathcal{X} , then f(x) satisfies the following.

Concavity: Let f be a real valued function over X, $f: X \to \Re$ is said to concave if $\forall x_1, x_2 \in X$ and $\lambda \in (0, 1), f((1 - \lambda) * x_1 + \lambda * x_2) \ge (1 - \lambda) * f(x_1) + \lambda * f(x_2)$. If f is twice differentiable and second derivative is non-positive over X then f is concave i.e. $f''(x) \le 0, \forall x \in X$.

Convexity: Let f be a real valued function over X, $f: X \to \Re$ is said to convex if $\forall x_1, x_2 \in X$ and $\lambda \in (0, 1), f((1 - \lambda) * x_1 + \lambda * x_2) \le (1 - \lambda) * f(x_1) + \lambda * f(x_2)$. If f is twice differentiable and second derivative is nonnegative over X then f is convex i.e $f''(x) \ge 0, \forall x \in X$.

Infection Point: The inflexion point of any real valued function over X, $f: X \to \Re$ is a point in X, at which the convexity switches to concavity vice versa.i.e the point $x \in X$ such that f''(x) = 0.

$$f(x) = \frac{1}{\sigma\sqrt{2\pi}} exp\left(\frac{-(x-\hat{\mathbf{w}})^2}{2\sigma^2}\right) \,\forall x \in \Re$$
$$\log(f(x)) = -\log(\sigma\sqrt{2\pi}) - \left(\frac{(x-\hat{\mathbf{w}})^2}{2\sigma^2}\right)$$

$$f'(x) = \frac{-(x - \hat{w})}{\sigma^2} f(x)$$
$$f''(x) = \frac{-\{(x - \hat{w})f'(x) + f(x)\}}{\sigma^2}$$

At the inflexion point f''(x) = 0, then we have

$$(x - \hat{\mathbf{w}})f'(x) = f(x)$$

$$(x - \hat{\mathbf{w}}) = \frac{f(x)}{f'(x)}$$

$$(x - \hat{\mathbf{w}})^2 = \sigma^2$$

$$x = \hat{\mathbf{w}} + \sigma$$

For the Fig.4, Mean and standard deviation of the weights are 2.665 and 1.265 respectively and hence the inflexion point of the MST is 2.665 + 1.265 = 3.9315. The edges which are greater than inflexion are pruned from MST. In the Fig.5, the values of the edges 4-5 and 1-5 are greater than 3.9315. So these edges are pruned from MST. This results into three disjoint sub-trees as shown in Fig. 6. Each sub-tree of forest represents as initial cluster (basic equivalence class) and the points in sub-tree are relatively close to each other.

Inducing the Equivalence classes: After splitting the MST forest into sub-trees, generalized equivalence class of each sub-tree are generated by using the attribute hierarchies. For example, in Fig 6, the QI values of sub-trees 1-2, 3-4 and 5-6-7 are {<21, M, 3051>, <26, M, 13053>}, {<33, F, 13063>, <37, F, 13068>}, {<45, M, 14851>, <48, F, 14856>, <50, M, 14865>} respectively and their generalized equivalence classes are <[20-30], M, 1305*>, <[30-40], F, 1306*>, <[40-50], Person, 148**>.

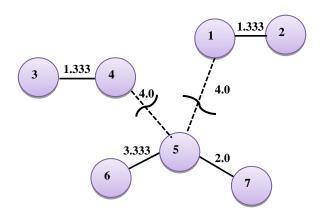


Fig 5: Formation of Sub-trees from MST forest

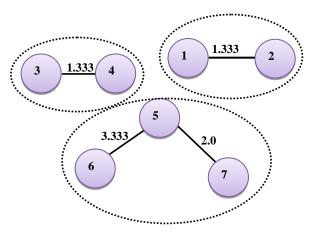


Fig 6: Initial Basic Clusters

5.5 Checking for *k*-anonymity

In this section we discuss the anonymity parameter checking of each equivalence class, which are generated previously. For example if the required anonymity constraint is k=2 for the Table 1(a), then it satisfied by the Table 1(b) formed by above step. That is each equivalence class contained at least two tuples which share the same value. If the anonymity constraint k=3 then Table 1(b) do not satisfied the anonymity constraint.

So that the equivalence classes, whose size is less than k are merged based on the following distance metric.

Definition 7 (Distance between two equivalence classes): Let C_1 be an equivalence class containing n_1 identical tuples t_1 and C_2 be an equivalence class containing n_2 identical tuples t_2 . t_{12} is the closest common generalization of t_1 and t_2 the distance between two equivalence classes is defined as follows

$$CDist(C_1, C_2) = n_1 * HDist(t_1, t_{12}) + n_2 * HDist(t_2, t_{12})$$
 (5)

We considered the anonymity requirement as k=3 for the table 1(a). The generalized table as shown in table 1(b) does not satisfy the anonymity constraint. Hence, we need to merge these equivalence classes based on the distance metric to satisfy the anonymity requirement.

Table 4. Definitions

CreateMST (QI)	Generates an MST from microdata
EdgeSet (T)	Returns the set of edges in minimum spanning tree T
AvgWeight (T)	Returns the average weight of edges in tree T
StdDev (T)	Return standard deviation of weights of edges in tree T
PartMST (T, Ē)	Returns set of sub-trees after removing the inconsistence edges \bar{E} from T .
GenClass(st)	Returns the generalized equivalence class of each sub-tree (cluster).

6. ALGORITHM

Our algorithm runs in two phases. During the first phase MST is constructed over the given microdata using Kruskal algorithm [21]. After constructing MST all the inconsistence edges are removed.

Algorithm 2: Minimum Spanning Tree based

k-anonymity(MSTK)

Input: Quasi Identifier tuples, Attribute Hierarchies and *k*

Output: Set of equivalence classes

Method:

- 1. Begin
- 2. $T \leftarrow CreateMST(QI)$
- 3. $E \leftarrow EdgeSet(T)$ // Set of edges of MST
- 4. $\bar{E} = \Phi$
- 5. $\hat{\mathbf{w}} \leftarrow AvgWeight(T)$

```
7.
         \sigma \leftarrow StdDev(T)
8.
        for each edge e \in E do
9.
             if (Weight (e) > \hat{\mathbf{w}} + \boldsymbol{\sigma}) then
10.
                  \bar{\mathbf{E}} \leftarrow \bar{\mathbf{E}} \cup \{e\}
11.
              end if
12.
         end for
          ST \leftarrow PartMST(T, \bar{E})
13.
             G \leftarrow \Phi
13..
14.
             G_1 \leftarrow \Phi
15.
             G_2 \leftarrow \Phi
16.
          for each sub-tree st \in ST do
17.
                  if |st| > k and |st| \le 2k - 1
                        G \leftarrow G \cup GenClass(st)
18.
                  else if (|st| > 2k - 1) do
19.
20.
                       call PartMST ( st, \bar{e} )
                  else if (|st| < k)
21.
                       G_1 = G_1 \cup GenClass(st)
22.
23.
           end for
           for each G_{1i}, G_{1j} in G_1 and i \neq j do
24.
25.
                 G_2 = G_2 \cup merge(G_{1i}, G_{1i})
                // G_{1i}, G_{1j} have minimum distance
                 G_1 = G_1 - \{G_{1i}, G_{1i}\}
26.
27.
           return (G + G_2) // return k-anonymized table.
29.
30.
         end Begin
```

In the second phase equivalence classes are generated and then the anonymity requirement is verified for each equivalence class. We defined various functions used in our algorithmic description of MSTK as shown in Table 4.

The quasi identifier point set (QI) and anonymity constraint k is the input to the algorithm. Initially we create MST tree (T) on quasi identifiers QI. The nodes of MST are QI tuples and weights are the distance between two tuples. After construction of MST, the tree T is partitioned by removing all edges whose weight w > AvgWeight(T) + StdDev(T). This results into set of sub-trees ST = {ST₁, ST₂,...} each representing an equivalence class. Let T' be the MST forest, obtained after removing inconsistent edges, i.e., $T' = \bigcup_i ST_i$ and $\bigcap_i ST_i = \emptyset$ where ST_i , $\forall i$ is a sub-tree of the forest (step 1-12). In the second phase, the k-anonymity in each equivalence class (sub-tree) is verified on the number of points in each sub-tree, Here three cases may arise.

Case 1: If the numbers of data points in the sub-tree lie between the k and 2k-1, then it satisfies the k-anonymity property(line 17-18). Now, GenClass(ST_i) is called to substitute appropriate domain ranges to all the data points of the subtree using attribute domain generalization hierarchies.

Case 2: If the number of data points in the sub-tree are greater than 2k-1, recursively we remove the inconsistent edges from the subtree until k-anonymity property is satisfied(line 19-20). We then generalize the data points by calling GenClass.

Case 3: If the number of data points in subtree are less than k, we generalize the data points in the subtree to form equivalence classes. In order to achieve k-anonymity, we merge the equivalence classes by calculating the pair wise distances between the equivalence classes (lines 24-27).

6.1 Complexity Analysis

The complexity for computing the distance matrix is $O(n^2)$ and to construct the MST from n number of data points is $n \log_2 n$. The maximum number of inconsistent edges that can be removed from MST forest are n-l, hence the time required for removing the inconsistent edges and to form sub-trees (equivalence class) is O(n). Let $C(C \ll n)$ be the total number of equivalence classes formed by after removing the inconsistent edges. Then the time complexity to check anonymity of each equivalence class is O(C). Therefore, the total time needed for executing the algorithm is $O(n^2) + O(n \log_2 n) + O(n) + O(n) + O(n) \approx O(n^2)$.

7. QUALITY MEASURES OF ANONYMIZATION

Protecting privacy of the data is achieved through anonymization. The other aspect of privacy is to produce the anonymized data that must be useful for deriving useful patterns for statistical analysis i.e., published data should remain for practical use. There are broad categories of information metrics for measuring the data usefulness. A data metric measures the data quality of the entire anonymous table with respect to the data quality in the original table. A quality metric guides each step of an anonymization (search) algorithm to identify an anonymous table with maximum utility or minimum information loss. However, some data publishers concentrate on anonymization process and some on utility of the data. To overcome this trade-off scenario the anonymized table is measured with reasonable information metrics to measure "similarity" between the original data and the anonymous data. This underpins the principle of minimal distortion. Several quality measure such as minimal distortion [3, 17], Information loss [26, 08], Discernibility metric [10], classification metric [27], information trade-off metric (entropy) [13], model accuracy and query quality [14] and normalized equivalence class metric [10] were widely used in the literature. These metrics are used based on anonymization problem. In this paper we used information loss, Discernibility and Normalized average equivalence class for measuring the quality of anonymize table.

7.1 Information Loss

Information loss is data metric proposed by Gabriel [8] to capture the information loss by generalizing a specific value to a generalized value. The information loss is computed based on attribute domain hierarchies. Information losses for both numerical and categorical attributes are given in equations 6 and 7 respectively.

$$IL_{numerical} = \frac{(U_i - L_i)}{(U - L)} \tag{6}$$

Here, $[U_i, L_i]$ be the upper and lower boundary values of the specific equivalence class and [U, L] be the upper and lower values of entire domain range satisfying the relation $[U_i, L_i] \subseteq [U, L]$.

$$IL_{categorical} = \frac{(L_v - 1)}{L} \tag{7}$$

Here, L_v is the number of nodes rooted from the current node and L is number of leaf nodes of the domain hierarchy. If the tuple consists of 'm' numerical attributes and 'n' categorical attributes then the information loss of the whole tuple is

$$IL_{tuple} = \sum_{i=1}^{m} IL_{(numerical)_i} + \sum_{i=1}^{n} IL_{(categorical)_i}$$
 (8)

The total information loss for the anonymized table T is

$$IL_{T} = \frac{\sum_{i=1}^{k} n_{i} * IL_{tuple}}{\sum_{i=1}^{k} n_{i}}$$
 (9)

7.2 Discernability

The discernibility metric (DM) [7, 10] addresses the notion of loss by charging a penalty to each record for being indistinguishable from other records with respect to quasi identifier group (equivalence class). If a record belongs to a quasi-identifier group of size |E|, then penalty for the record will be |E|. The overall penalty cost of generalized table T is given by

$$DM = \sum_{Equivalence \ classes \ E} |E|^2 \tag{10}$$

The objective of anonymization is to minimizing discernibility cost. For example, in the anonymized table of our running example Table 1(b) contain three equivalence classes {<[20-30], M, 1305*>, <[30-40], F, 1306*>, <[40-50], Person, 148**>} and their group sizes are 2, 2, 3 respectively. The discernibility of the table is $2^2 + 2^2 + 3^2 = 17$.

7.3 Normalized Average Equivalence Class

The intuition of this metric is to measure how well the partitioning approaches the best case when each tuple is generalized in a group of k indistinguishable tuples. i.e.,

$$CAVG = \frac{total\ number\ of\ records}{total\ number\ of\ equivalence\ classes\ *\ k}\ (11)$$

The objective of anonymization is to reduce the normalized average equivalence class size. If CAVG is low, more number of equivalence classes will be generated and hence we can achieve better anonymization.

8. EXPERIMENTAL RESULTS

A battery of experiments on Intel core 2 duo @ 2.93 GHz with 1.5GB RAM being allocated separately for the Netbeans platform was run. The benchmark dataset, Adult dataset from the UC Irvine Machine Learning Repository was used [24] for

experimentation. The size of the dataset was 30,162 tuples after removing tuples with missing values

	Attribute	Туре	Distinct values	Tree height
1	Age	Numeric	74	4
2	Work class	Categorical	7	3
3	Gender	Categorical	2	1
4	Education	Categorical	16	4
5	Race	Categorical	5	2
6	Occupation	Categorical	14	2

We considered the projection of the Adult dataset with six attributes Age, Work Class, Gender, Education, Race and Occupation. The corresponding tree heights of the chosen attributes and their type and number of distinct values are shown in the Table 3.

The proposed method is compared with a global recoding method basic Incognito [20]. Comparison is done based on different quality metrics like information loss, discernability metric and normalised average equivalence class size. Experiments are shown in Fig 6(a), Fig 6(b) and Fig 6(c).

Fig 6(a) shows the information loss when the value of *k*-increases for both incognito and our approach. Our method shows a significant improvement by reducing the information loss when compared to incognito.

Fig 6(b) shows how discernability metric (DM) differs for both incognito and MSTK. For DM value, out of 7 experiments (k=3...21) with step size of k=3, MSTK outperforms Incognito. We found that, our approach achieves better discernability.

Fig 6(c) shows that our approach achieves better normalized average equivalence class than incognito. For lower value of k the incognito's Normalised average equivalence class size (CAVG) is very high where as in our minimum spanning tree based k-anonymity (MSTK) approach the CAVG is quite low for any value of k. The running time complexity however is more for our approach but a promising quality is achieved when compared with incognito. We can decrease the time by using parallel algorithms for construct the minimum spanning tree.

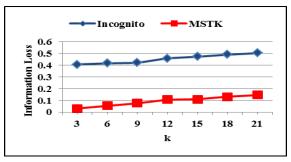


Fig 6(a): Information Loss

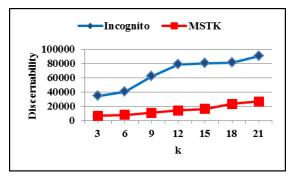


Fig 6(b): Discernability

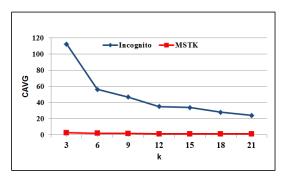


Fig 6(c): Normalized Average Equivalence Class

Fig 6: Performance of Different Methods with Variant k

9. CONCLUSION AND FUTURE WORK

In this paper, we studied the problem of minimum spanning tree based technique to achieve k-anonymity. We compared our approach with incognito technique based on quality metrics like information loss, CAVG and discernability. Our technique showed a significant improvement when compared with incognito. The complexity of our algorithm is O (n^2). The experimentation clearly signifies the quality that can be achieved using our approach. In future we would like to enhance our MSTK method by adopting parallelism to construct the MST and thereby reducing the running time.

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